Oscar, Brayden, Nathan

Math 251, Fri 20-Nov-2020 -- Fri 20-Nov-2020 Discrete Mathematics Fall 2020

RSA encryption needs · fast modular exponentiation . solving modular equations

Friday, November 20th 2020

Wk 12, Fr

Topic:: Euclidean algorithm Read:: Rosen 4.3 Topic:: Solving congruences Read:: Rosen 4.4 Arithmetic mod m (need integer $m \ge 2$) Define \mathbb{Z}_m : both a set of integers $\{0, 1, 2, 3, ..., m-1\}$ along with operations $a + b = a + b \mod m$ $a \cdot b = ab \mod m$

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$$27 \cdot 4(-5) = 3 \cdot 4(-1) = -3 \text{ mel} 4$$
$$= 1$$

In R, the solution x to ax = 1
is called the multiplicative inverse of a.
In mod 4 arithmetic, 3 is its own multiplicative
inverse.

Solving congruence equations

Ex. I Solve
$$3x + 1 \equiv 2 \pmod{4}$$

 $3x \equiv 1 \pmod{4}$ ofter subtracting (
NGW, Since 3 is multiplicative inv. (mod 4)
of itself, can multiply both sides by 3
 $3 \cdot 3x \equiv 3 \cdot 1 \pmod{4}$
 $5 \sqrt{x} \equiv 3 (\mod{4})$
 $5 \sqrt{x} \equiv 3 (\mod{4})$
 $(\operatorname{really means all of these solve:
 $\cdots = 5 - 1, 3, 7, 11, 15, \cdots$$

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$$19 \times = 51 \pmod{420}$$

Underlying question: What is 19's mult. inv. (mod 420)?

Euclidean Algorithm

Our main purpose in learning it: You use it to determine gcd(a, b). How it works:

1. Start with two (usually positive) integers *a*, *b*. Call the larger one r_0 , the smaller r_1 .

- 2. Iterate the division algorithm:
 - Divide r_0 by r_1 —that is, find integers q_1 and r_2 (Note: $0 \le r_2 < r_1$) so that

$$r_0 = q_1 r_1 + r_2$$

Shift roles *r*₁ into the former role of *r*₀, *r*₂ into the former role of *r*₁, and repeat, using the division algorithm to find *q*₂ and *r*₃. (Note: 0 ≤ *r*₃ < *r*₂.)

$$r_1 = q_2 r_2 + r_3$$

We continue to shift roles of the r_j and repeat. This process produces a strictly decreasing sequence

$$r_0, r_1, r_2, \ldots, r_n$$

until a remainder, call it r_{n+1} , finally is zero, which is our **stopping criterion**.

Example 1:

Perform the Euclidean algorithm with a = 276, b = 324, to find gcd(276, 324). Start by labeling larger number r_0 , smaller r_1 . $r_0 = 324$, $r_1 = 276$. $r_0 = 7, -8$ $324 = 276 \cdot 1 + 48$ Notes: Any divisor of two-out-of-three in list r_0, r_1, r_2 is a divisor of all three. $r_2 \leq r_1$

See website https://www.extendedeuclideanalgorithm.com/calculator.php.



Example 2:

Perform the Euclidean algorithm with a = 4312, b = 585.

$$r_{0} = 4312, \quad r_{1} = 585$$

$$4312 = \frac{7}{2}, (585) + \frac{217}{217}$$

$$585 = \frac{2}{2}, (217) + \frac{157}{3}$$

$$217 = \frac{1}{(151)} + \frac{66}{66} - \frac{5}{4}$$

$$151 = \frac{2}{2}, (66) + \frac{19}{4} - \frac{5}{3}$$

$$66 = \frac{3}{2}, (19) + \frac{9}{4} - \frac{5}{4}$$

$$19 = \frac{2}{4}, (9) + \frac{1}{4} - \frac{5}{4}$$

$$9 = \frac{9}{4}, (1) + \frac{9}{4}$$

Extended Euclidean Algorithm

Our main purpose in learning it. to write gcd(a, b) as a linear combination ta + bKy theorem based on our Eucl. alg. work EX.] Above we found gcd(324, 276) = 12In proparation, silve for remainder $48 = 36 \cdot 1 + 12 \rightarrow 12 = 48 - 36$ $276 = 48 \cdot 5 + 36 - 36 = 276 - (48)5)$ $324 = 276 \cdot 1 + 48 - 548 = 324 - 276$ Now carry out extended Enc. Alg - Writing $12 = t \cdot 324 + b \cdot 276$ $w/ b, t \in \mathbb{Z}$.

$$12 = 48 - 36$$

$$= 48 - [176 - (48)(5)]$$

$$= (48)(6) - 276$$

$$= (324 - 276)(6) - 276$$

$$= (324)(6 + (276)(-7))$$

$$7 + (276)(-7)$$

$$7 + (5 + 6) = -7$$